Descriptive Set Theory
Lecture 26

Classitication problems.

Lit $X$ be a collection of arathenaticd objects (seg, Rienam iartacos, mearce-pres automorphisa, (*.algebras,..) and let $E$ be an equivaleces celation on $X$ (e.g, isomorplisen). A (lansification probtion is to understand the objects in $X$ up to the equir ael. E. Ileall, we would like bo have a "reasoncble" assigmat

$h: X \rightarrow \mathbb{R}$ lor sone other nice space, e.g. $Z^{\mathbb{N}}$ )


$$
x_{1} E x_{2} \Leftrightarrow h\left(x_{1}\right)=h\left(x_{2}\right) .
$$

Objects arisimg in cualy $h_{c}$ subjects, sulte as analysis, differential geonety, harmonic annalssis, dyuanies (neasunable, topological, imooth), opecator algehress ( $C^{*}$, von Nenwanne), and furctional analysis, can often be encoded into a Polish space, i.e. X is Polife. For exaple, we already law tht $K\left(\left[0,17^{1 N}\right)\right.$ and $F\left(\mathbb{R}^{\mathbb{N}}\right)$ can be thought of as the Polishe rpaces of coupnect sefarable ipaces $d$ of Polish spaces, resp.

Morcover, the en. ael. E that we care about we unst otkn analytic (as sabuets of $\mathrm{X}^{2}$ ) become thay we defied wing an existential quantifyier over a Polish followed by southing Boed. As for the assignouent $h$, one can alwags use axiem of boice to pick a point troo each dars al yet a desieced functien to $\mathbb{R}$ by cardisulith cossideratious. Deananding $h$ to be untinoons is too anch leze then building 4 may be obsitcucted b, he top on $X$ dich has wothin, to do vith the anplexity of the closification poblem, i.e. he couplexits of $E$. $S_{\text {? }}$ we demand $h$ to le Bored.

Def. An ey. re. E on a stond. Boal pare $X$ is called concectily damitiable (ar smooth) if $\partial$ Borel rechuction $h$ of $E$ to equality $=\mathbb{R}$ on $\mathbb{R}$ (esuiv. any othe uncthl st. Boul space), i.e. $\exists L: x \rightarrow \mathbb{R}, t . \forall x_{x}, x_{2} \in t$,

$$
x_{1} E_{x_{2}} \Leftrightarrow h\left(x_{1}\right)=h\left(x_{2}\right) .
$$

I. oter vords, $h$ lesendes to an enbedding $x / E \in \mathbb{R}$.

More yenecally, solving a clanitication problen $E$ or $X$ weans to cumerstand the "Bonel cardinalif" of $X / E$, i.. we say $|X / E| S_{B}|Y / F|$, whee $B$ stonds for Brel, if $\partial$ Bond unap $~ h: X \rightarrow Y$ that lescends $t$ an ingetion $X / E \leftrightarrow Y / F$. Is paticalar, for $E$ on $X$, $X / E$ ung larger Burcl curclinalits than $X$.

Def. We suy hat an uy. nel. E on a st. Boel $X$ is Bovel reducible to an eq. nel. Fon a st. Bor. Y, if $\exists$ Borel recluction $h$ of $E$ do $F$, ie. $\exists L: X \rightarrow Y$, t. $\forall x_{1}, x_{2}+X$,

$$
x_{1} E_{x_{2}} \Leftrightarrow h\left(x_{1}\right) F h\left(x_{2}\right) .
$$

In oher vords, $h$ descends to an eubedding $X / E \leftrightarrow Y / F$. We denote this by $E \leqslant B F$.

The study of Boal rechacibilitg of analytic as. rel. has becone its own subject the provides colonitication and anti-clansification zssilts to the aforementioned vicas of watth. This happened in the last 30 gears, initiated 4) Alexander Kechcis, A. Lonvean, Grey Hjorth, and othes.

Note that if $E$ on $X$ is suooth then $E$ is Bosel (as a subsit of $X^{2}$ ): $E=\tilde{h}^{-1}\left(\Delta_{\mathbb{R}}\right)$, shere $\Delta_{\mathbb{R}}=$ diagoanal al $\tilde{h}\left(x_{1}, x_{2}\right):=\left(h\left(x_{1}\right), h\left(x_{2}\right)\right)$, so $\tilde{h}$ is Bovel.
This "wens"s" the the is a cithl algorith i.t. given $x_{1}, x_{2} \in X$, detecmises chethe $x_{1} E x_{2}$. In fact, re can tatee a sequence of yes/ao Bonel questions s.t. $x_{1} E x_{2}$ iff $x_{1}$ al $x_{2}$ 's ansuers to Vase gacstions are the sare:

Prop. $E$ on $X$ is snoth $\Leftrightarrow \nexists$ Bouel sets (guestions) $X_{a_{i}^{c}} Q_{n} \subseteq X$ s.t. $\forall x_{1}, x_{2} \in X$,


$$
x_{1} E x_{2} \Leftrightarrow \forall_{n}\left(x_{1} \in Q_{n} \Leftrightarrow x_{2} \in Q_{n}\right)
$$

$Q_{0}^{c}$ Proof. $<$. We define a Bonel reduction of $E$ $a_{0}{ }^{c}$
to $=z^{N}$ by: $x \mapsto\left(a_{n}^{x}\right)_{n}$, here
$a_{n}:=\left\{\begin{array}{ll}0 & \text { if } x \notin Q_{n} \\ 1 & \text { if } x \in Q_{n}\end{array}\right.$ Incleed,
then $\left(a_{n}^{x_{1}}\right)_{n}=\left(a_{n}^{x_{2}}\right)_{n} \Leftrightarrow x_{1} E x_{2}$.
$\Rightarrow$ Inppose $a$ Boal coducton $h: X \rightarrow 2^{N}$ of $E t_{0}=z^{(N)}$. Define $Q_{n}:=h^{-1}\left(\left[+*+\ldots+\frac{1}{h}+\ldots\right]\right)$.

Examples of smooth eq. el.
(a) Isomorphism of fire gen. abel- yes. Firstly, we anode all (tl gps with the underling set $\mathbb{N}$ into a Polish space as follows: a group $P$ or $\mathbb{N}$ is a stenchare $(\mathbb{N}, \cdot)$ where $\cdot$ is a binary op, satisfying so e axioms. Replying . by its graph, ie. a ternary ablation $R \leq \mathbb{N}^{3}$, we get a relational structure $(\mathbb{N}, R$.$) , so P$ can be decoded frow R.. R. $\in 2^{\left(\mathbb{N}^{3}\right)}$ al those R. $\in 2^{\left(N^{3}\right)}$ WA satisfy the group cxions ten a losel ut. Thus, $\sigma_{p}(\mathbb{N})$ is a coupact Polish space. The spacelvat fir. gen. abel sp form a $\sum_{3}^{0}$ subset of $C_{p}(\mathbb{N})$, hence it's a st. Boer l space. We know foo- algebra the eve 3 fin, gees. ch. Sp $P$ is som. Do a sp of the form $\mathbb{Z}^{n} \times($ tina .aha. sp $)$. Taurus at the the -ap $\Gamma \mapsto(n$, finabgp $)$ is, Bowel, cituening the smoothness of the isom. rel. on FCA. ${ }^{\left(n^{2}\right)}$
(b) Let $M_{n}(\mathbb{C})$ be the (Polish) space of all $n \times 5$ couple -atrices. kt $\sim$ denote the similarity of matrices, ie.
matrices $A \sim B \Leftrightarrow A, B$ ane corjegate

$$
\Leftrightarrow f^{\prime} Q \in G L_{n}(C) Q A Q^{-1}=B
$$

By bef, $\sim$ is an wanalitic en. eel.
Ltting $J(A)$ denote the Jordan caronical hor of $A$, me huen trom lin. als. Wht $A \sim B \Leftrightarrow J(A)=J(B)$. Again, one can show $U_{t}$ the wap $A \mapsto S(A)$ is Bovel, vittressing the snoothaess (hence also (Boulmon) of $\sim$.

Prop. Let $E$ be an eq. ael on a Polich $X$. It $E$ is gencrically eis. (ic. E-inv. Booll sths we neasee ar coreager) and cach $E$-clan is wreager, then $E$ is nonsunoth. Sinilarly, if $E$ is r-agodic I ench $E$-dam is $\mu$-uall, for soe Borel measane $\mu$ on $K$, then $E$ is noninooth.
Prod. To pore both at one, call meager/null uets small.
Kprose $\in S$ shooth, so $\exists$ Bonel $h: X \rightarrow 2^{\mathbb{N}}$ s.t. $\forall x_{1}, x_{2} \in X_{1}$

$$
x_{1} E x_{2} \Longleftrightarrow h\left(x_{1}\right)=h\left(x_{2}\right),
$$

i.c. L is contant in each E-clon, in particularer, $h^{-1}(B)$ is Einvaciant Bonel, for eanh $B \leq 2^{N}$. Thas, $h^{-1}(B)$ is suall or cosmall. Call an $s \in 2^{<N}$ heavy if $h^{-1}([s])$ is cosuall. If $s$ is heavy then $\exists$ (unigue) $i \in\{0,1\}$ s.t. si is heary. Start with $\varnothing$, which is hereug al wase it lown to $x \in 2^{N}$, follxing the wavy Rild. But $\{x\}=\bigcap_{4}\left[\left.x\right|_{n}\right\}$ so $h^{-1}(x)=\bigcap_{n} h^{-1}\left(\left\{\left.x\right|_{L}\right\}\right)$ is cosmall, hat $h^{-1}(x)$ is at most one E-chen, Mich is small, a coutcacliction.

Exngles: $E_{\mathbb{Q}}(\mathbb{Q} \gtrdot \mathbb{R}$ by tionslation), ireational cotution ane uonimoolh. So is $\mathbb{E}_{0}$ on $2^{i N}$ defined $b_{s}$ $x \mathbb{E}_{0} y \Leftrightarrow \forall_{n}^{\infty} \quad x(n)=s(n) \quad$ (eventad agalith). Note tht if $\mathbb{E}_{0} \leq_{B} E$, then $E$ is also won ruooth (othernise, copposition would wituen swoothues of $\mathbb{E}_{0}$ ). Turss out, this is the only obstiction to snootheres for Boul ey rel.
$\mathbb{E}_{0}$-dichotomg (Kechris-Harricyton-loweana) For an Bocel ey. rel. $E$ on a $i t$. Bonel $X$ either $E$ is suooth, i.e. $E \leq_{B}=2^{N}$. or $\quad \mathbb{E}_{0} \subseteq_{B} E$.


This showe tht $\mathbb{E}_{0}$ is the minimun element anong nonsmooth Bonel ey. sel. This dich gevecalites earlier Thans of Glima al Effros, ro it is afercel as the generclized Glimm-Effros diclutong.
the original proof of this dichatons ases effective esscriptive theorg, i.e. a tiver top. on $X$ that womes tron Turing unchines.

